

Part 2

Part 2 contains:

- Some theory
- A hands on explanation

Please ask questions!

Important Words

- A Monomial is a finite product of unknowns.

$$xy^5z^3$$

- A Polynomial is a finite sum of monomials.

$$4x^5 + 3x^3 + 1$$
$$xyz^5 + 7x^2y + 1$$

Ideal

The ideal of $\{f_i\}$ is the set

$$I(\{f_i\}) = \left\{ \sum_{\text{finite}} f_j p_j \mid f_j \in \{f_i\}, p_j \in \mathbb{C}[x] \right\}$$

- Contains all polynomials we can generate.
- All the polynomials in the ideal are zero on our zero set.
- Tells us about the number of solutions.
- Contains an infinite number of polynomials.
- If the number of solutions is finite there is a monovariate polynomial in this set.

Polynomial Division and rests

- For univariate division is well defined
- For multivariate division the order of monomials is a central issue.

$$\frac{y^2 + x}{x + y} = 1 + \frac{y^2 - y}{x + y} = y + \frac{-xy + x}{x + y}$$

Is $x > y^2$?

Monomial Orders

Polynomial division requires that the monomials have an order. There exist several monomial orders

- Lexicographic order (phone-book order), very intuitive, gives long computations. Sort by power of x , then by power of y .
- Graded Lexicographic order. Sort first by total degree and then by lexicographic order. More intuitive and quite good for computing Gröbner bases.
- Graded reversed Lexicographic order, very counterintuitive but it is easier to compute a GB.

Gröbner Basis (GB)

- A basis for the ideal.
- Depends on the monomial order chosen.
- Gives the number of solutions.
- There exist algorithms to compute it.
- Gives ways to compute solutions if the number of solutions is finite.
- All leading terms of polynomials of the ideal are divisible by the leading term of some polynomial in the GB.

S-Polynomials

Definition: Given two polynomials f_1 and f_2 with leading terms $x^{\alpha(1)}$ and $x^{\alpha(2)}$. Let x^β be the least common multiple of the leading terms, then

$$S(f_1, f_2) = \frac{x^\beta}{x^{\alpha(1)}} f_1 - \frac{x^\beta}{x^{\alpha(2)}} f_2$$

is the S-polynomial of f_1 and f_2 .

- Multiply both polynomials so that their leading terms cancel out.
- Generating S-polynomials is a key step to building a Gröbner basis.

$$S(x^2 + 1, xy + 3) = y(x^2 + 1) - x(xy + 3) = -3x + y$$

Computing a Gröbner basis

There are many algorithms but the main parts are

- Remove linear dependency.
- Remove multiples of known polynomials.
- Generate new S-polynomials.
- Terminate when nothing new can be generated.
- Must use exact arithmetic.
 - Takes time
 - Takes memory

Matrix Representation of polynomials

$$\begin{aligned} &2x^2 + 4xy + y^2 + 8x + y + 1 \\ &2x^2 + 3xy + 2y^2 + 4x + 4y + 4 \end{aligned}$$

can be represented as

$$\begin{bmatrix} 2 & 4 & 1 & 8 & 1 & 1 \\ 2 & 3 & 2 & 4 & 4 & 4 \end{bmatrix} \begin{bmatrix} x^2 \\ xy \\ y^2 \\ x \\ y \\ 1 \end{bmatrix}$$

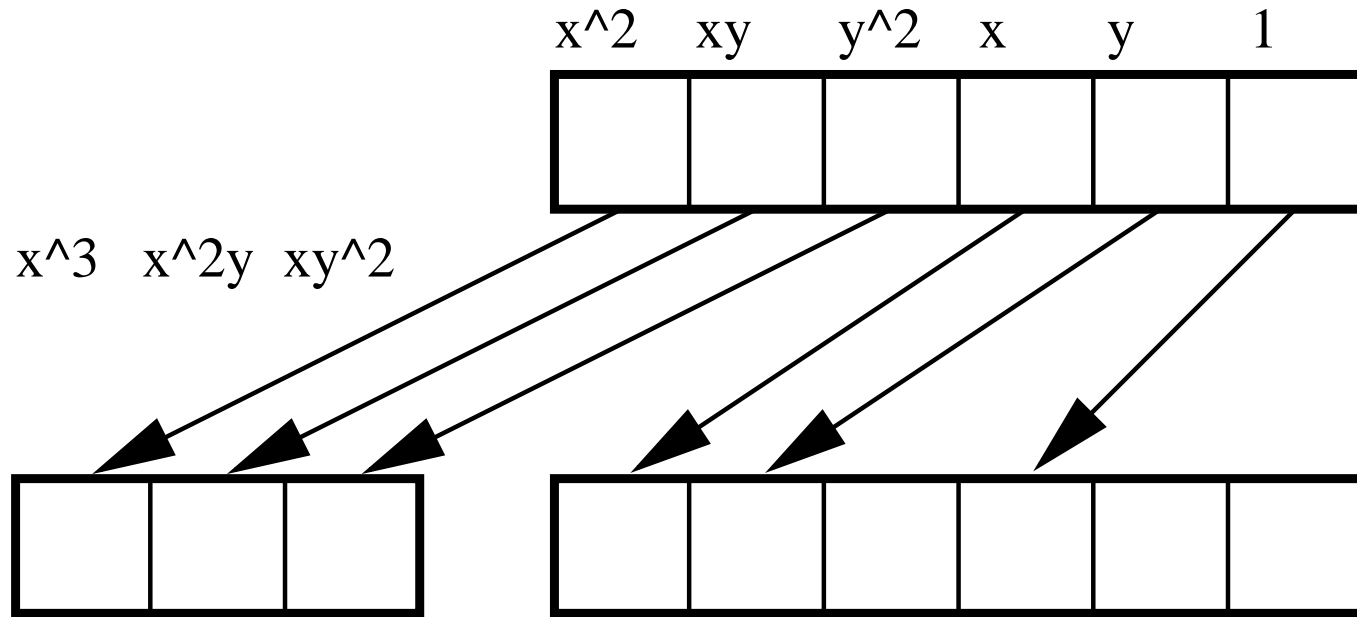
Multiplication with a monomial

- Multiplication is a shift operation

$$(4x + 5y + 6)x = 4x^2 + 5xy + 6x$$

$$\begin{bmatrix} 0 & 0 & 0 & 4 & 5 & 6 \end{bmatrix} \begin{bmatrix} x^2 \\ xy \\ y^2 \\ x \\ y \\ 1 \end{bmatrix} \longrightarrow \begin{bmatrix} 4 & 5 & 0 & 6 & 0 & 0 \end{bmatrix} \begin{bmatrix} x^2 \\ xy \\ y^2 \\ x \\ y \\ 1 \end{bmatrix}$$

Multiplication with a monomial



- Multiplication is a linear operator.

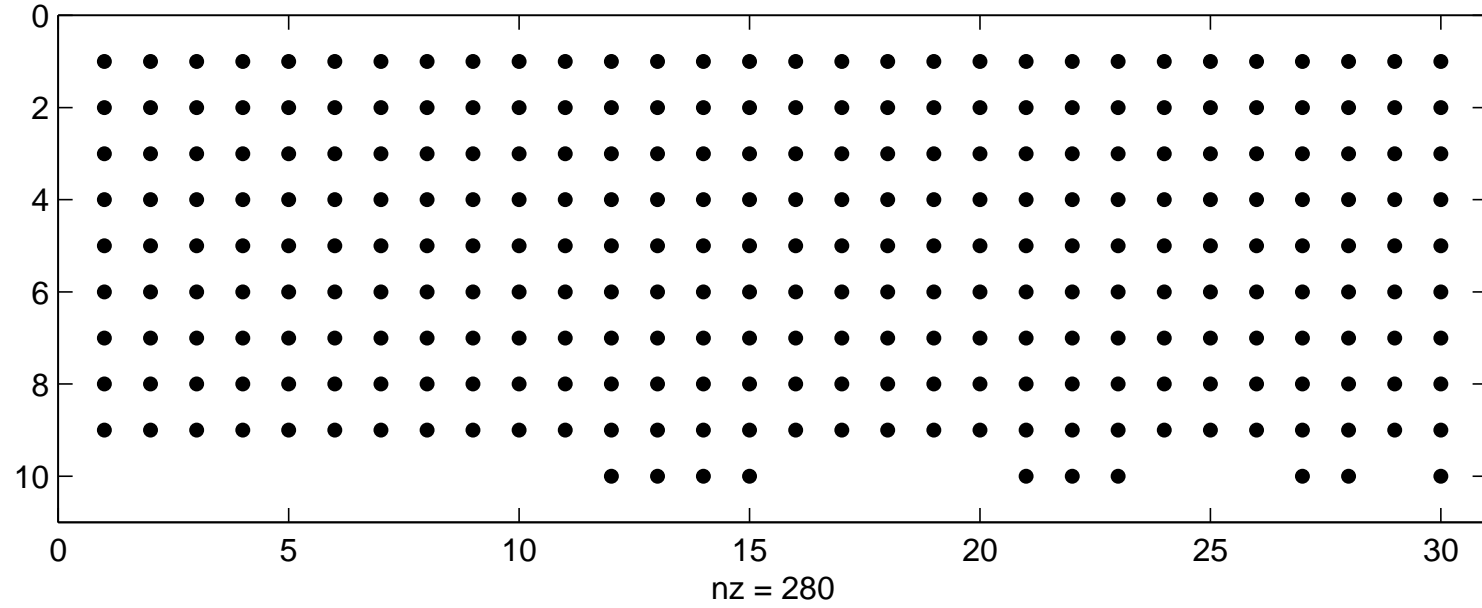
Matrix Based GB Computations

- Add lots of polynomials every time
- Use linear elimination to remove dependencies

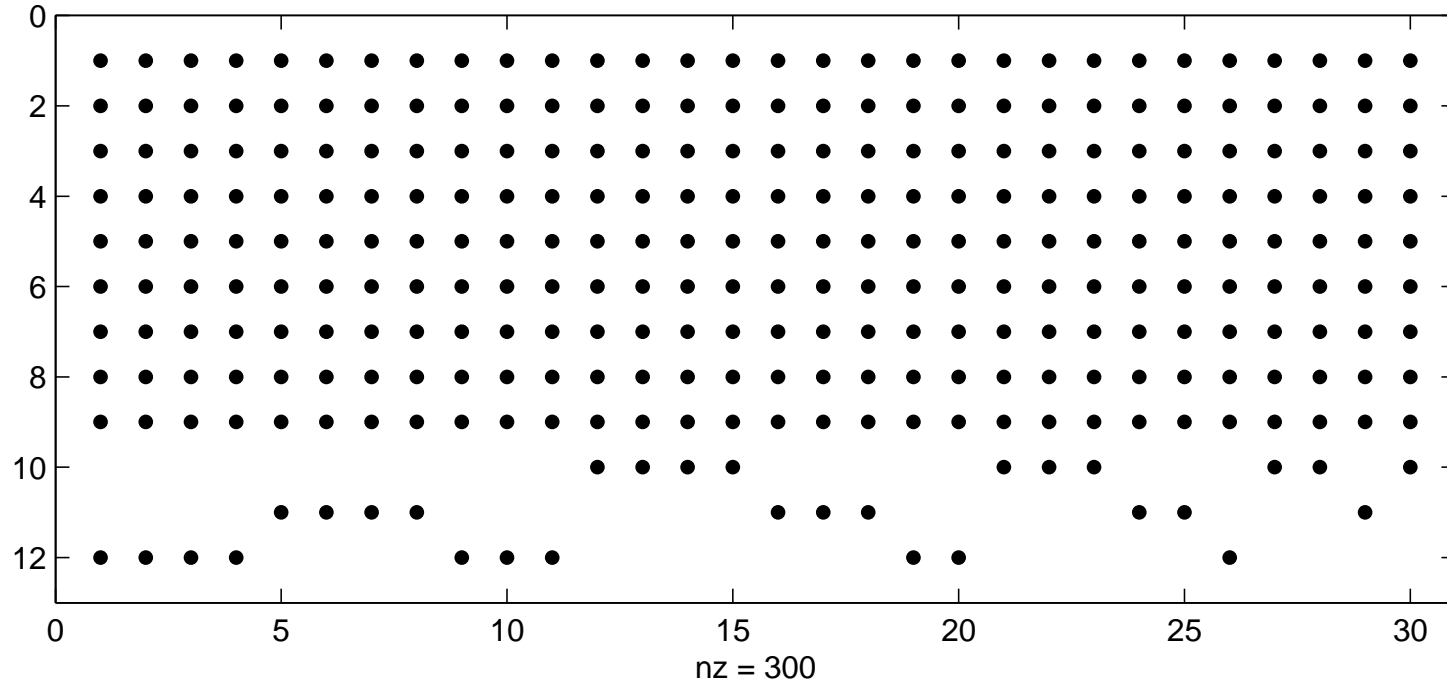
Faugere (F4), Rouillier . . .

- Example: six points and two cameras, unknown focal length.

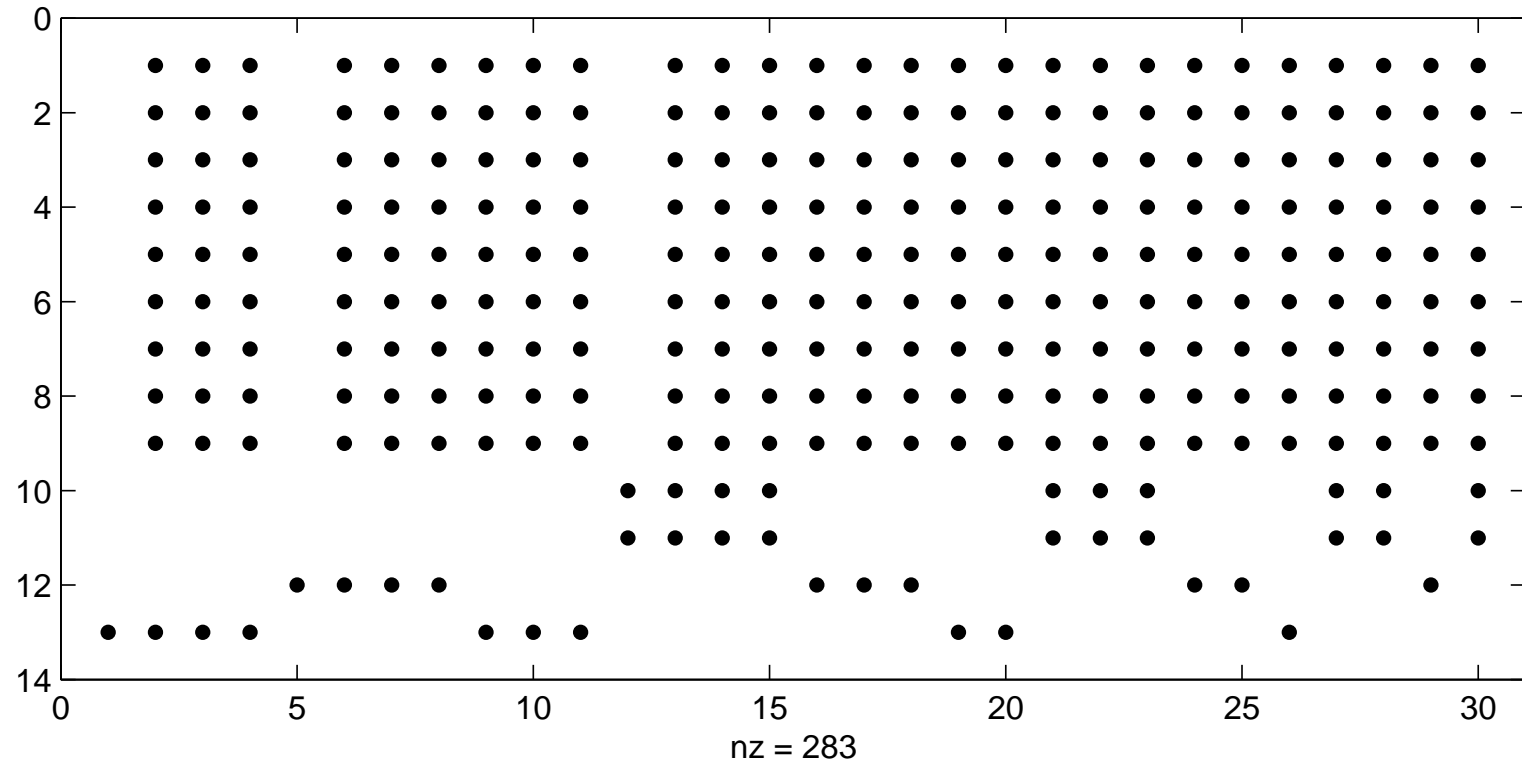
Computational Step



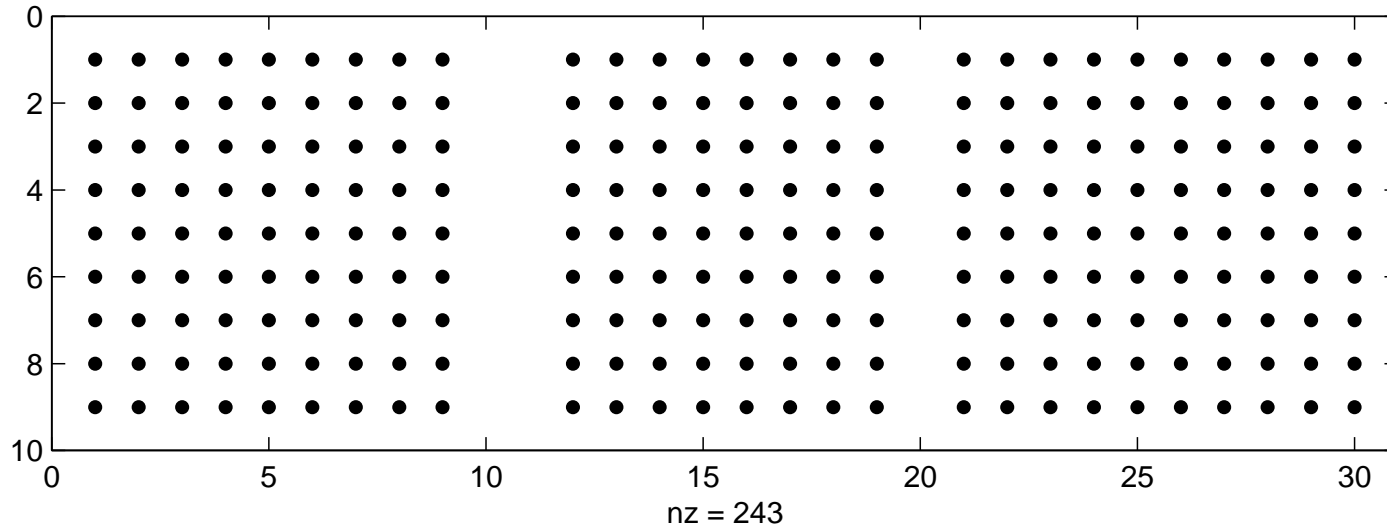
Computational Step



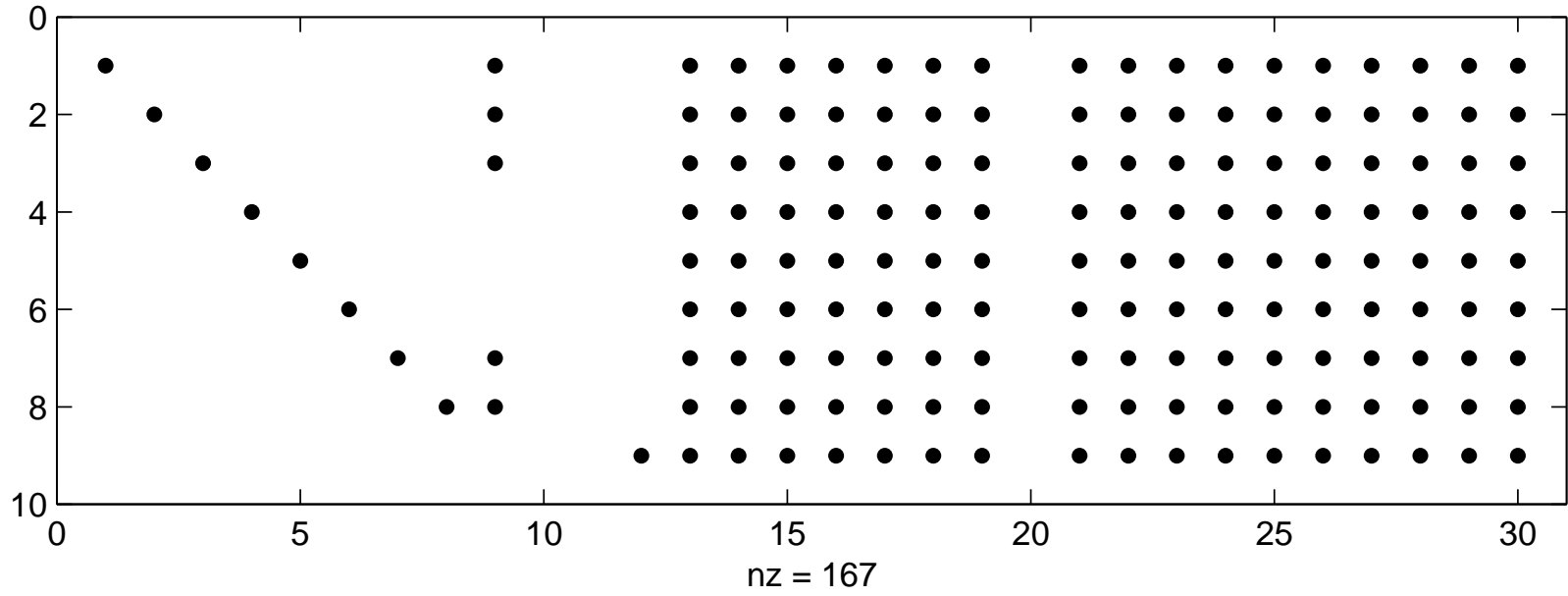
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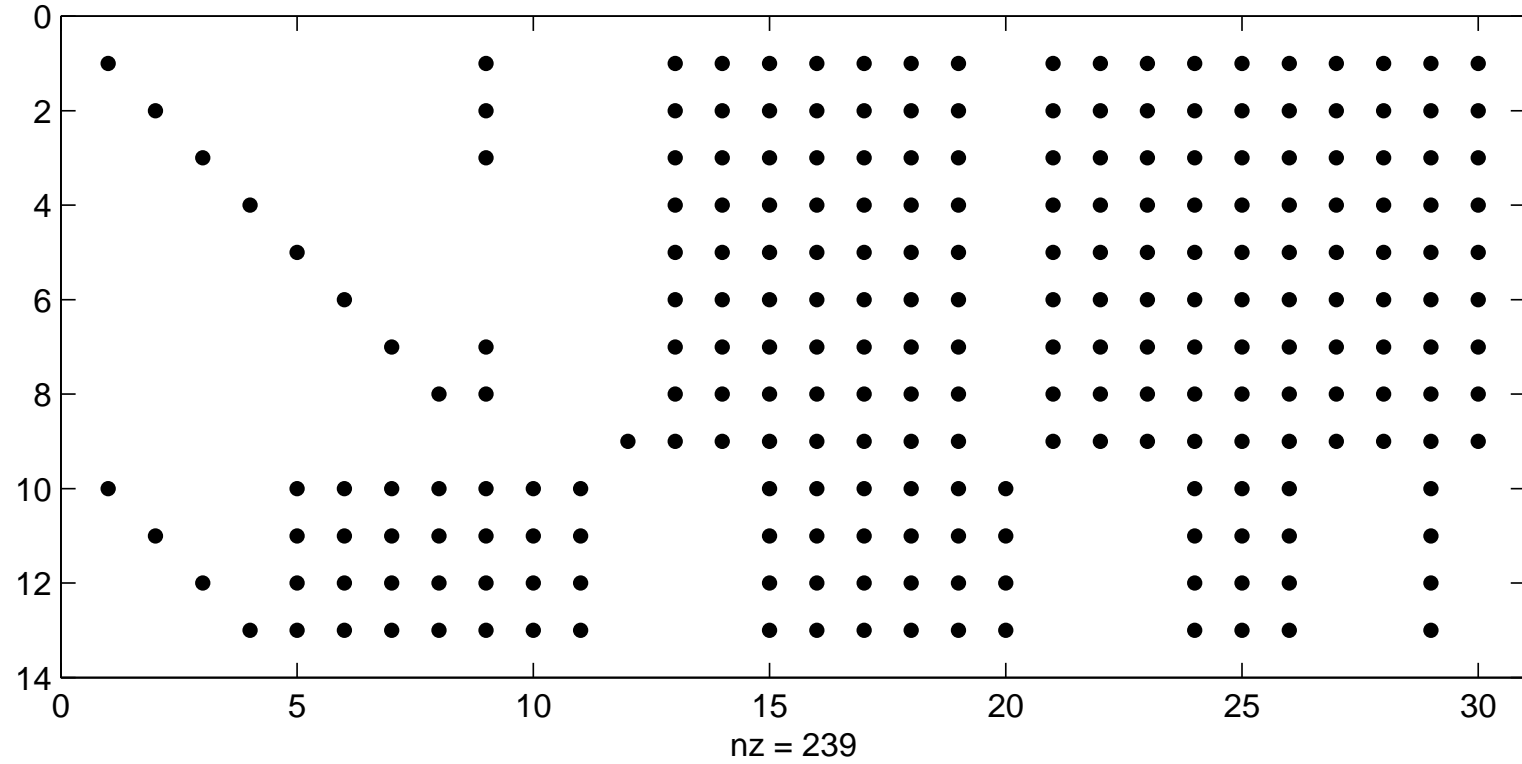
Computational Step



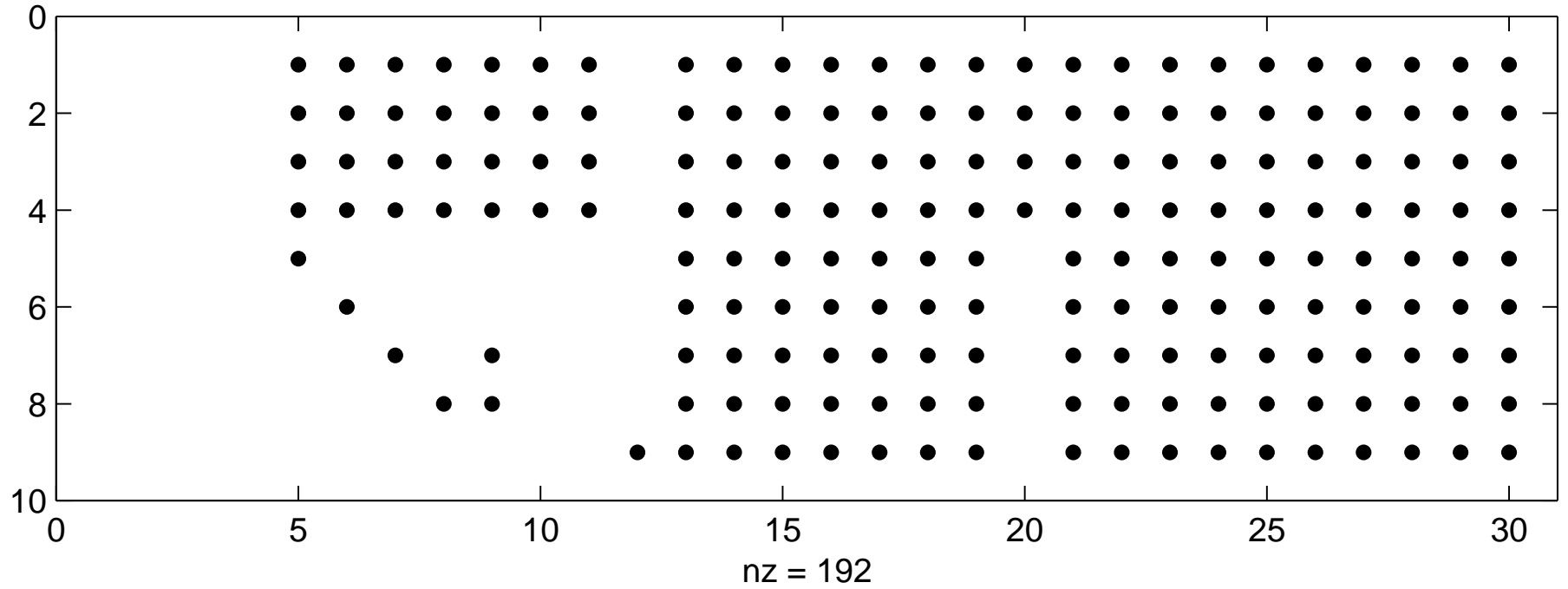
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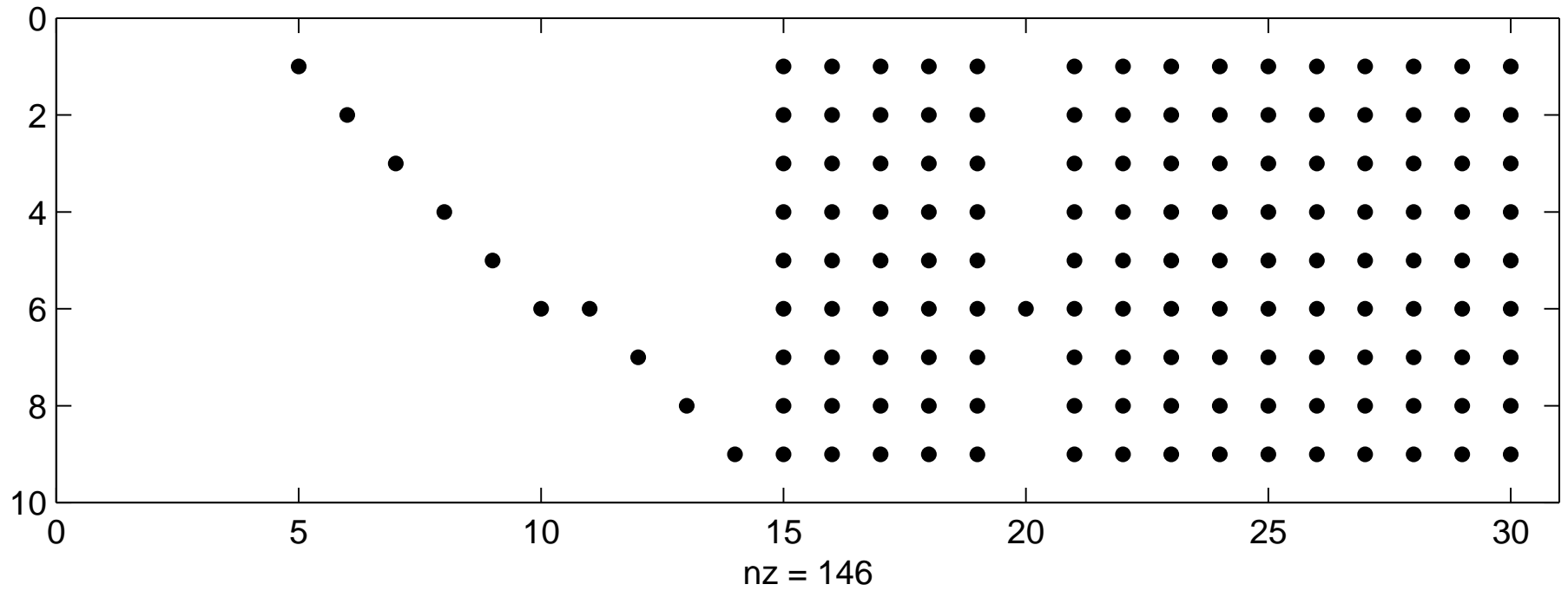
Computational Step



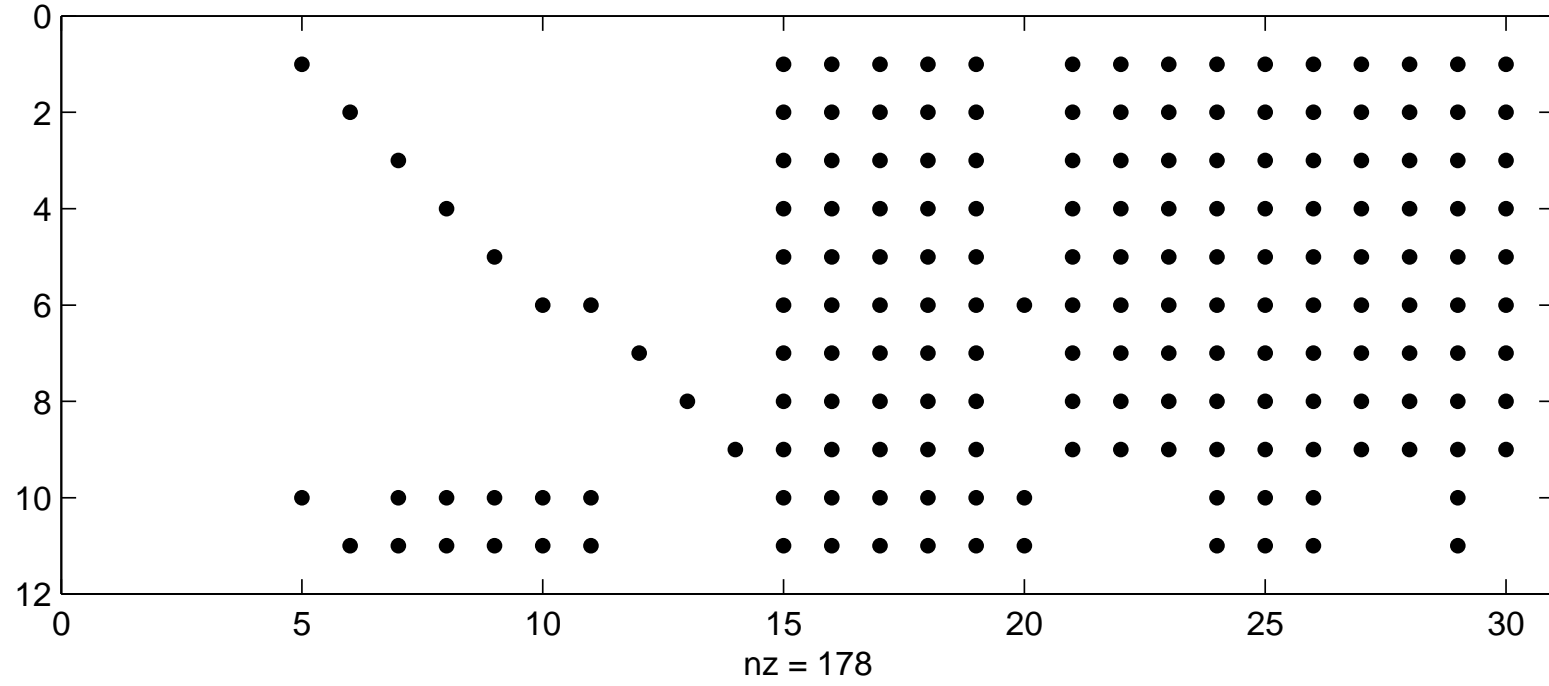
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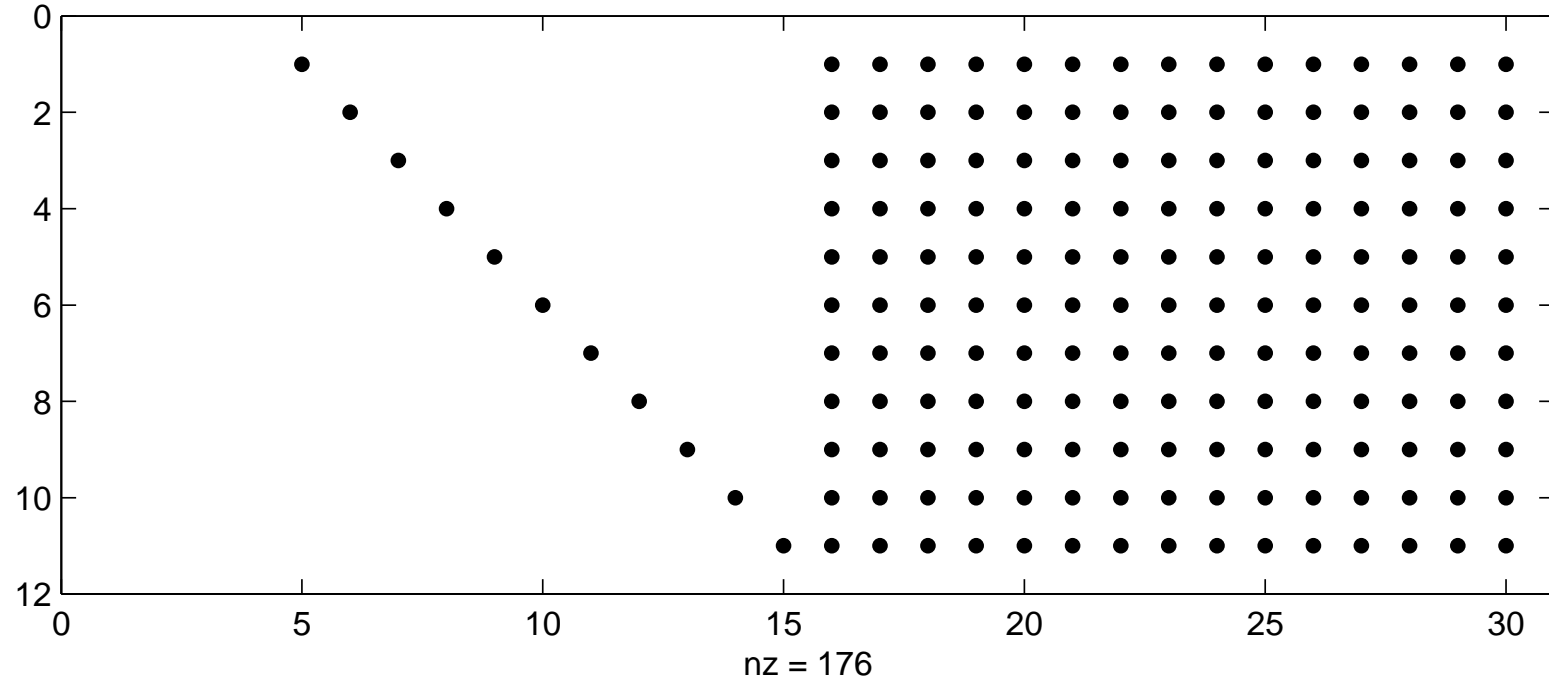
Computational Step



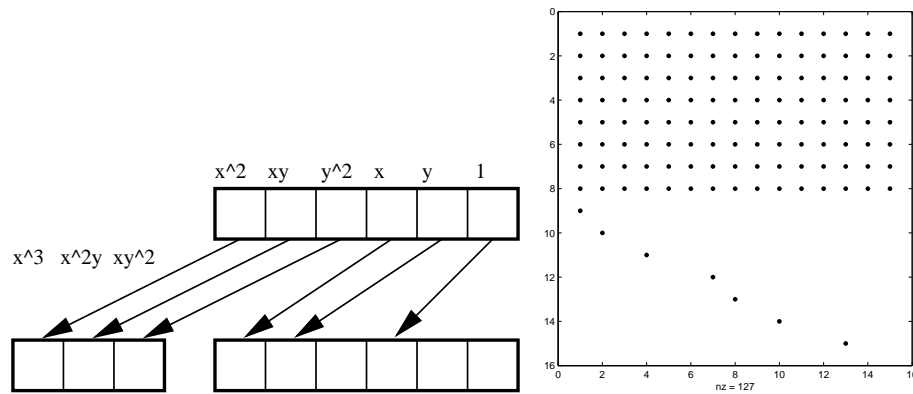
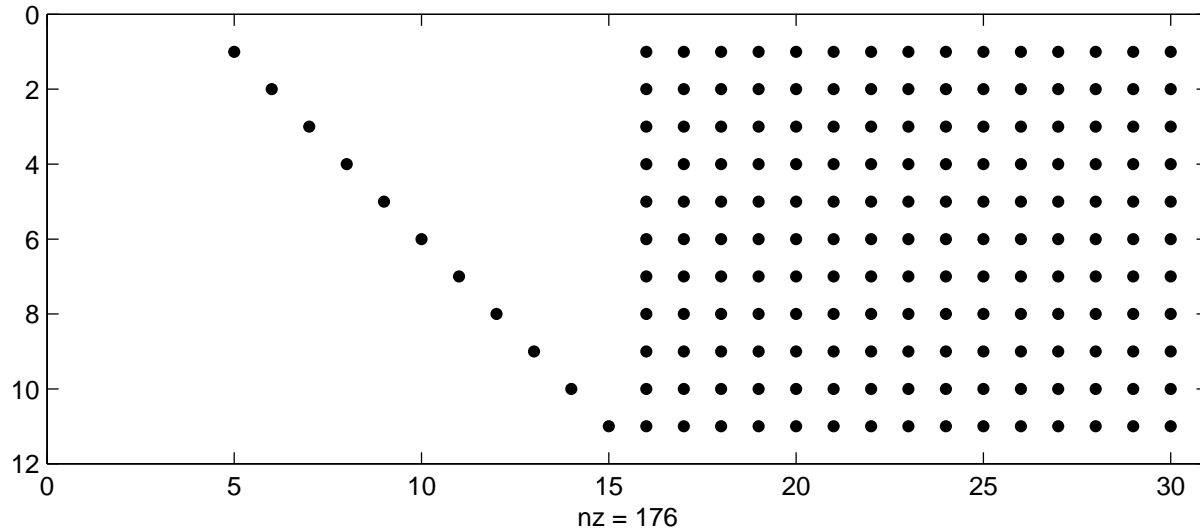
Computational Step



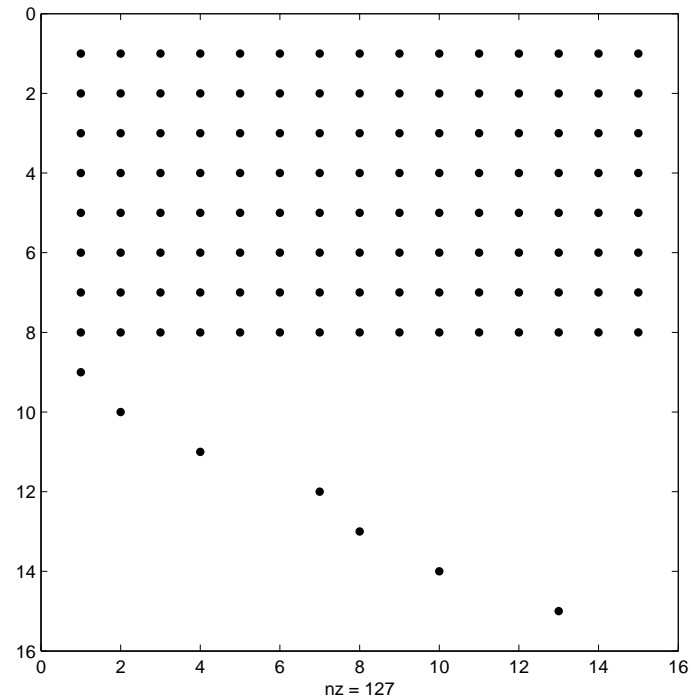
Computational Step



Extracting the Action Matrix



The Action Matrix



- The eigenvectors of this system gives the 15 solutions

A finite Field : \mathbb{Z}_5

The integers modulo a prime p is a finite field. (Modulo is the rest with division by p .)

- Closed under $+$ and \cdot
- Additive and multiplicative inverses

Addition table

	0	1	2	3	4
0	0	1	2	3	4
1	1	2	3	4	5
2	2	3	4	5	6
3	3	4	5	6	0
4	4	5	6	0	1

Multiplication table

	0	1	2	3	4
0	0	0	0	0	0
1	0	1	2	3	4
2	0	2	4	1	3
3	0	3	1	4	2
4	0	4	3	2	1

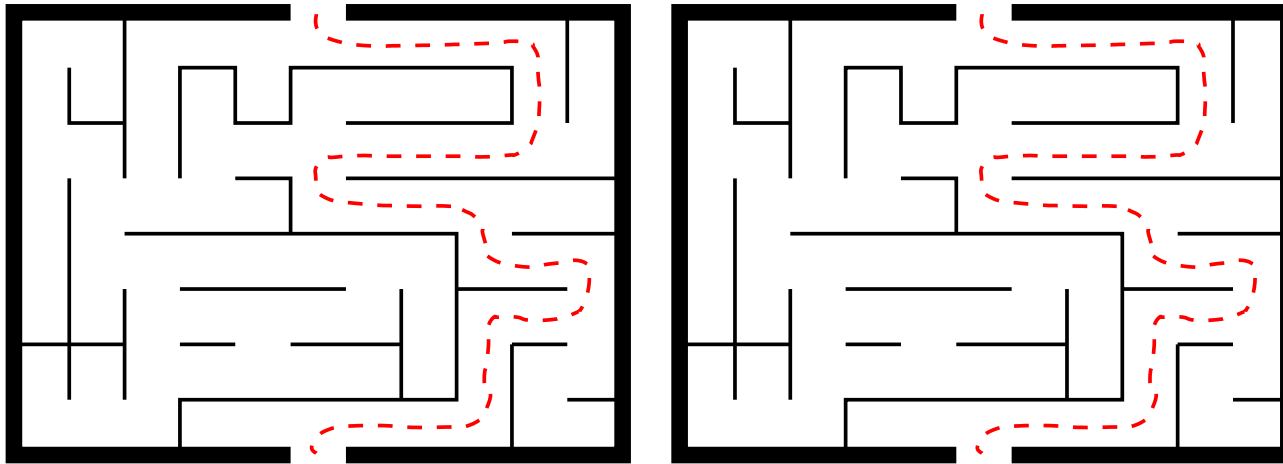
Why use a finite field?

- Exact arithmetic
 - No problem determining if an element is 0
- Constant memory size for all elements of the field
 - It takes more time to run out of memory
- Gives an excuse to brush of some memories of elementary number theory

Fields are usually similar

\mathbb{R}

\mathbb{Z}_p



- $p = 30097$ is a nice prime.
- Using \mathbb{Z}_p makes it easy to compute the trace.

The Quotient Ideal

- $\mathbb{C}[x]/I$, that is the equivalence classes modulo the ideal I .
 1. Take a polynomial in $\mathbb{C}[x]$
 2. Keep computing rests with polynomials in I until this is no longer possible.
 3. The rest is in $\mathbb{C}[x]/I$
- Example $\mathbb{C}[x]/(x^2 - 1)$
 - 2-dimensional space
 - A basis is 1 and x

The action matrix

Consider what happens to the basis polynomials of the quotient ideal $\mathbb{C}[x]/I$ under multiplication by a polynomial f . This is a linear operator A_f .

- Example: Consider $\mathbb{C}[x]/(x^2 - 1)$ with $f(x) = x$. Then $1 \rightarrow x$ and $x \rightarrow 1$.

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

- Note the eigenvalues! 1 and -1 .

Solution Procedure

1. Pose Problem
2. Count variables and constraints
3. Find GB-trace
4. Compute GB
5. Extract the action matrix
6. Solve Eigenproblem

Helpful tools

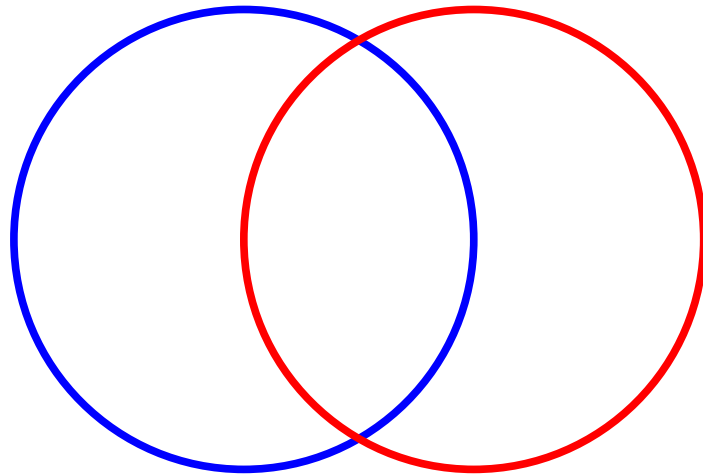
- Macaulay 2,
<http://www.math.uiuc.edu/Macaulay2/>
Daniel R. Grayson and Michael E. Stillman, provided
under GPL.
- perl
- matlab

Probably all the tools can be substituted for similar tools.

Finding Solutions

- Find all solutions to

$$\begin{cases} x^2 + y^2 = 1 \\ (x - 1)^2 + y^2 = 1 \end{cases}$$



Howto build an analogous problem 1

```
R = ZZ/30097[x,y];  
I = ideal( {x^2+y^2-4 , (x-1)^2+y^2-4 } );  
dim I , degree I
```

- We get 0,2

Looking at the trace

```
gbTrace 3
```

```
I = ideal( {x^2+y^2-4 , (x-1)^2+y^2-4 } );
```

```
dim I , degree I
```

```
Gives
```

```
m m m r
```

```
o12 = (0, 2)
```

m means an S-polynomial which did not reduce to 0.

r reduced to 0

o reduced to 0

Looking at the trace

```
gbTrace 3
```

```
I = ideal( {x^2+y^2-4 , (x-1)^2+y^2-4 } );
```

```
dim I , degree I
```

```
Gives (filtered a little)
```

```
--- computing pair (. . x2) ----
```

```
.0 gb 0 = x2<0>+y2<0>-4<0>
```

```
--- computing pair (. . x2) ----
```

```
gb 1 = x<0>+15048<0>
```

```
--- computing pair (1 0 x 1 x2) ----
```

```
gb 2 = y2<0>-7528<0>
```

```
--- computing pair (2 1 x y2 xy2 marked) ----
```

```
Number of pairs = 4
```

```
Number of gb elements = 2
```

Implementing this trace

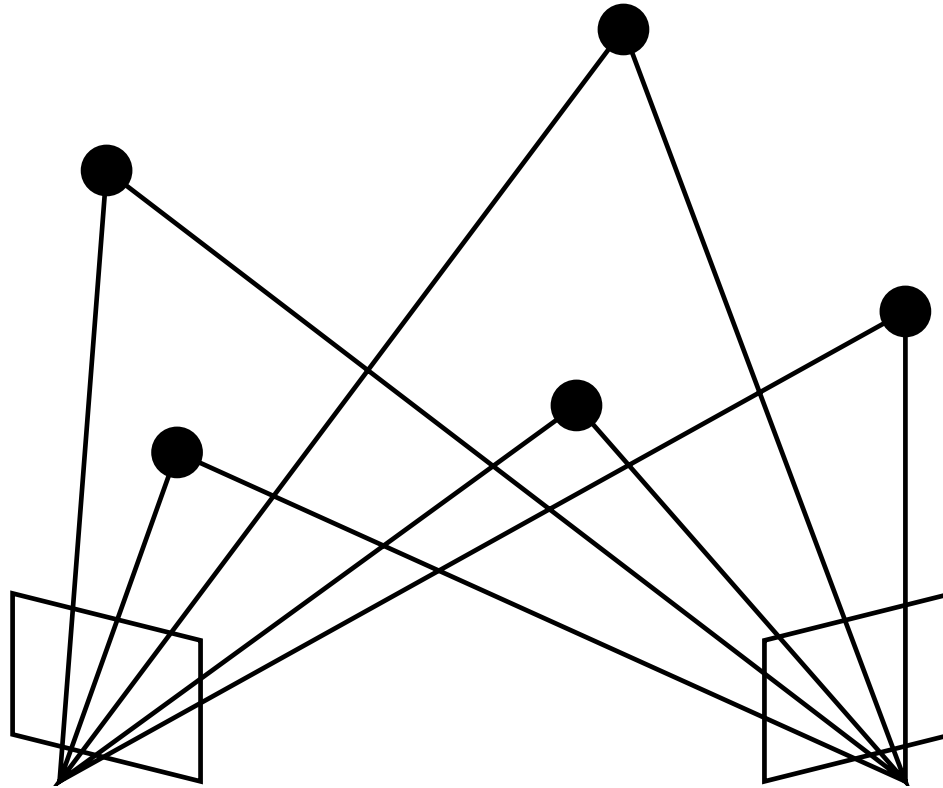
1. Start from $x^2 + y^2 - 1$ and $(x - 1)^2 + y^2 - 1$.
2. Linear elimination

$$\begin{bmatrix} 1 & 1 & 0 & -1 \\ 1 & 1 & -2 & 0 \end{bmatrix} \begin{bmatrix} x^2 \\ y^2 \\ x \\ 1 \end{bmatrix} \implies \begin{bmatrix} 1 & 1 & 0 & -1 \\ 0 & 0 & 1 & -0.5 \end{bmatrix} \begin{bmatrix} x^2 \\ y^2 \\ x \\ 1 \end{bmatrix} \implies \begin{cases} x^2 + y^2 & -1 \\ & x - 0.5 \end{cases}$$

3. Add $x(x - 0.5) = x^2 - 0.5x$ to the basis
4. Linear elimination

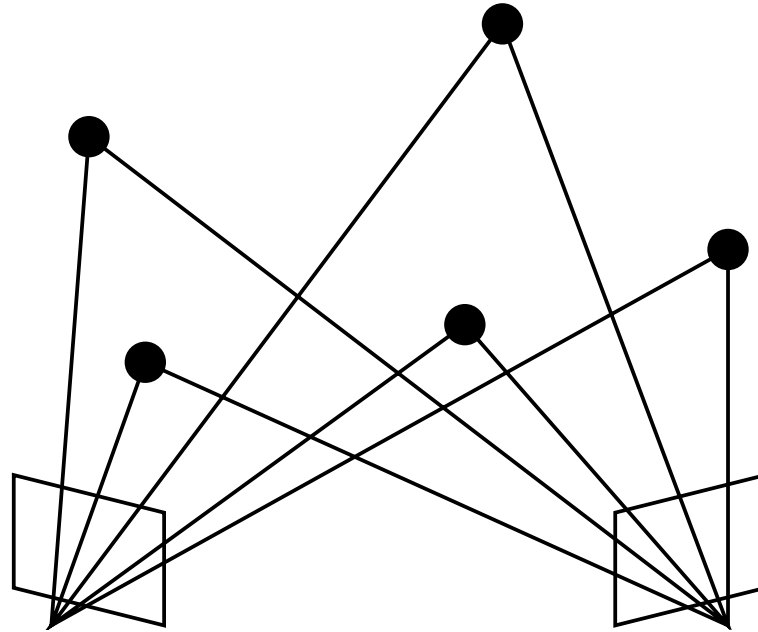
$$\begin{bmatrix} 1 & 1 & 0 & -1 \\ 0 & 0 & 1 & -0.5 \\ 1 & 0 & -0.5 & 0 \end{bmatrix} \begin{bmatrix} x^2 \\ y^2 \\ x \\ 1 \end{bmatrix} \implies \begin{bmatrix} 1 & 0 & 0 & 0.25 \\ 0 & 1 & 0 & -0.75 \\ 0 & 0 & 1 & -0.5 \end{bmatrix} \begin{bmatrix} x^2 \\ y^2 \\ x \\ 1 \end{bmatrix} \implies \begin{cases} x^2 & +0.25 \\ & y^2 - 0.75 \\ & & x - 0.5 \end{cases}$$

Relative Pose



- Observation Equation $\lambda_{ij}u_{ij} = P_i U_j$.
- Alternative Observation equation $x^\top F x' = 0$.
- Restrictions on P or F .

Finding Solutions



- Five points and two cameras relative pose

$$\begin{cases} E = E_0 + xE_1 + yE_2 + zE_3 \\ 2EE^T E - \text{trace}(EE^t)E = 0 \\ \det(E) = 0 \end{cases}$$

Howto build an analogous problem 2

Investigating the analogous problem for the 5-point problem.

```
R=ZZ/30097[x_0..x_2];
```

```
t={1_R,x_0,x_1,x_2};
```

```
E=fold((a,b)->a+b,apply(4,i->t_i*random(R^3,R^3)));
```

```
EEt = E*transpose(E);
```

```
I = ideal( 2*EEt*E-trace(EEt)*E )+ ideal det E;
```

```
gbTrace 3
```

```
dim I , degree I
```

Howto build an analogous problem

- Set up the finite field (3 variables)

$$R = \mathbb{Z}/30097[x_0 \dots x_2];$$

- Declare a list

$$t = \{1_R, x_0, x_1, x_2\};$$

- Set $E = E_3 + \sum_{i=0}^2 t_i E_3$ where E_i random.

$$E = \text{fold}((a, b) \rightarrow a + b, \text{apply}(4, i \rightarrow t_i * \text{random}(R^3, R^3)))$$

- Write down the equations, take the ideal

$$EEt = E * \text{transpose}(E);$$

$$I = \text{ideal}(2 * EEt * E - \text{trace}(EEt) * E) + \text{ideal det } E;$$

- Set the amount of information from the computation

$$\text{gbTrace } 3$$

- Compute the dimension and degree of the ideal

$$\text{dim } I, \text{ degree } I$$

More output

Set gbTrace 100, rerun computations and filter

```
--- computing pair (. . x_0^3) ----  
.0 gb 0 = x_0^3  
.  
.  
--- computing pair (. . x_0^3) ----  
gb 9 = x_2^3  
--- computing pair (9 8 x_1 x_2 x_1x_2^3) ----  
.  
.  
.21--- computing pair (1 0 x_0 x_1 x_0^3x_1) ----  
Number of gb elements      = 10  
Number of pairs computed   = 25
```

We see that no new S-polynomials were generated.

Implementing this trace

1. Set the 10 equations into a system with coefficients ordered in GrLex order.

```
MONS = [3 2 2 1 1 1 0 0 0 0 2 1 1 0 0 0 1 0 0 0;  
        0 1 0 2 1 0 3 2 1 0 0 1 0 2 1 0 0 1 0 0;  
        0 0 1 0 1 2 0 1 2 3 0 0 1 0 1 2 0 0 1 0];
```

2. Perform Gaussian elimination

$$\begin{bmatrix} I & M \end{bmatrix}$$

3. We now have a Gröbner basis

To get a complete solver we have to extract the action matrix.

<http://www.vis.uky.edu/~stewe/FIVEPOINT/>

The filter I used

```
rm tt
cat class25_new.m2 | M2 > tt
cat tt|perl -p -e 's/gb/##gb/;s/comput/##comput/;s/\<0\>.*//'
      |egrep "##" | perl -p -e 's/##//' | egrep -v marked
```

(one line-break inserted where there should be none)

Conclusions from M2

- Study the number of solutions
- Study if the trace looks possible
- This is an iterative process

The Problem Formulation is Very Important

- Get as many equations as possible
- Keep the degree as low as possible
- Keep the number of variables as low as possible

Tricks

- Use fundamental/essential matrix if possible.
- Use determinants to remove variables.
- ...

6 point, unknown f

- Stewénius, Nistér, Kahl and Shaffalitzky, CVPR2005.
- Using the trace constraint and the rank constraint building a GB is very straight forward.
- 15 solutions

6 point, unknown f

$$F = F_0 + l_1 F_1 + l_2 F_2$$

```
KK = ZZ / 30097;  
R = KK[l_1, l_2, p];  
F0 = random(R^3, R^3);  
F1 = random(R^3, R^3);  
F2 = random(R^3, R^3);  
F = F0 + l_1 * F1 + l_2 * F2;
```

6 point, unknown f

$$E = K^T F K$$

```
P = matrix {{1,0,0},{0,1,0},{0,0,p}};  
tr = X -> det( X_{0}^{0} +X_{1}^{1}+X_{2}^{2});  
EQ1=2*F*P*transpose(F)*P*F-tr(P*F*P*transpose(F))*F;  
EQ2=det( F);  
I = ideal EQ1 + ideal EQ2;  
gbTrace 3;  
dim I , degree I
```


Implementing a Trace

We have a formulation where the trace looks ok

1. Write a matrix version in M2.
2. Build the equations in real numbers
3. Port the matrix version to Matlab or C

A Matrix Version in M2

- M2 was not meant for matrix operations
- Need to hack our own Gaussian elimination
 - It can be done using the Gröbner engine
 - It is also possible to hack it using special commands. Please email me for source.

Extracting the equations

Two alternatives

- Symbolically expand the equations and identify the coefficients (slower)
- Build a vector based multiplication, much like convolution. (faster)

It is strongly recommended to use some symbolic toolbox for this. For more complex expressions the expansion can be very long!

Porting a solver

1. Put the original equations into a matrix
2. Derive the multiplication matrices
3. Work carefully. check against a problem where the ground truth is known.

Setting up the monomials in GrevLex order

```
function MONS = build_MONS_3var(n)
MONS = [];
for i1=0:n
    for i2=0:n
        if i1+i2 <= n
            MONS = [MONS [i1;i2]];
        end
    end
end
end
[dummy,I]=sortrows( [sum(MONS); -MONS(end:-1:1,:) ]
MONS = MONS(:,I(end:-1:1));
```

Build a multiplication matrix

```
function M = build_mult_matrix( MONS_from,MONS_to,mult)
% MONS_* matrices where the columns represent monomial
% mult is a column matrix representing a monomial
%
mult = mult(:)';
M = [];
for i=1:size(MONS_from,2)
    p = MONS_from(:,i)';
    M=[M ismember( MONS_to', p+mult , 'rows' ) ];
end
M = M';
```

Extract the Action Matrix

For the 5-point the action matrix is read by the following commands.

```
M = -A([1 2 3 5 6 8], :);
```

```
M(7,1) = 1;
```

```
M(8,2) = 1;
```

```
M(9,4) = 1;
```

```
M(10,7) = 1;
```

- Use code to build this code! (does not fit this page)

Action Matrix to Solutions

```
[V,D] = eig(M );  
SOLS = V(7:9,:) ./ (ones(3,1)*V(10,:));  
and then back-substitute  
Evec = EE*[SOLS ; ones(1,10) ];
```

planarly restricted generalized camera

- H. Stewénus and K. Åstrom ECCV2004, before we fully understood how to work with GB.
- There are two easy minimal relative pose problems.
- This case is suitable as an easy exercise to begin with.

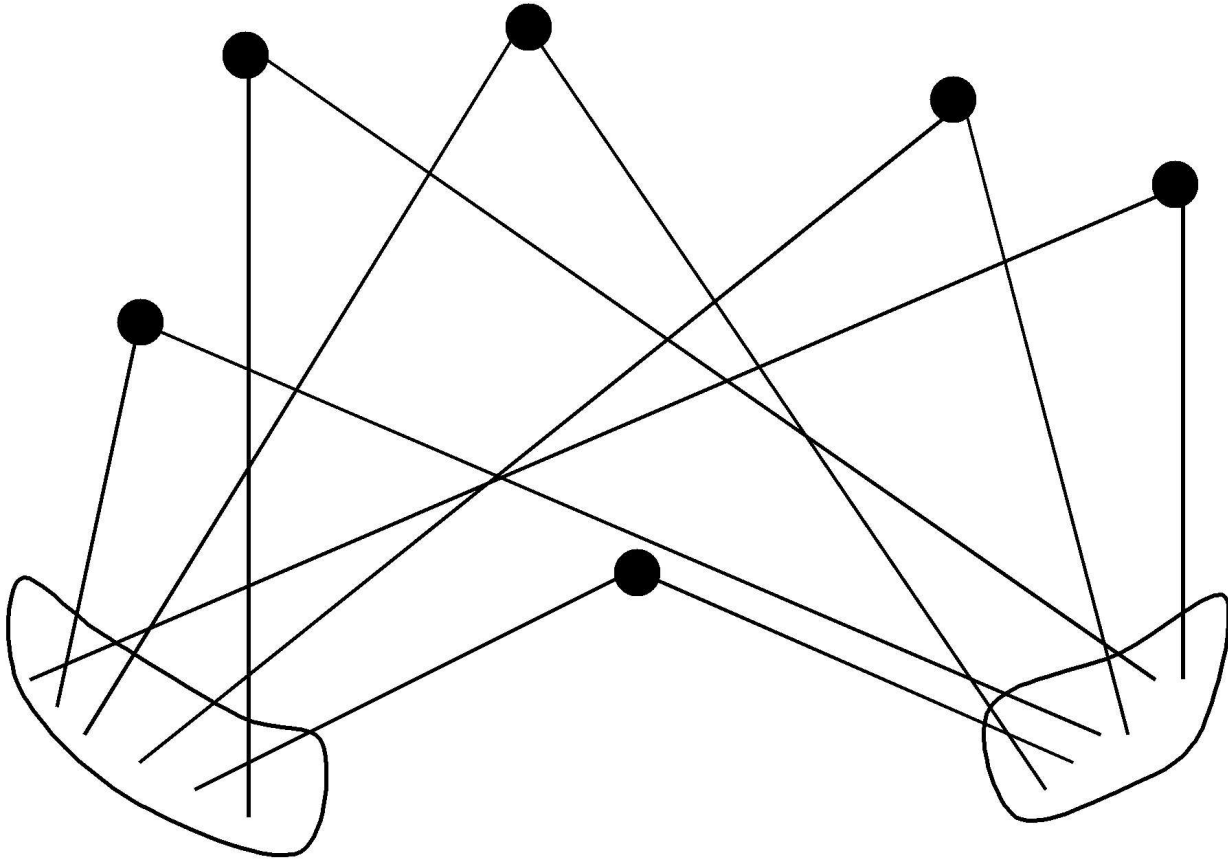
Relative Pose generalized camera

- There are 2 minimal cases, 64 solutions and 1320 solutions.
- For the easy problem:
 - Stewénius, Nistér, Oskarsson and Åström at OMNIVIS 2005.
 - Departing from the generalized epipolar constraint.
 - Requires a rather hard-to-find rewrite

The Generalized Camera

- A Generalized Camera maps from points in space to lines in the camera coordinate system.

$$(X, T) \rightarrow (q, q')$$



The Generalized Epipolar Constraint

- $R_1 = I$
- R_2 is the rotation at time i
- t_i is the translation at time i
- (q_i, q'_i) is the observation in the camera

$$q_2^\top R_2 q'_1 + q_2^\top (R_2 [t_1]_\times - [t_2]_\times R_2) q_1 + q_1^\top R_2^\top q'_2 = 0$$

- This gives a polynomial

Parameterizing with depth i

Let $U_1 = \begin{bmatrix} 0 & 0 & 0 & 1 \end{bmatrix}^T$ define the coordinate system. Then

$$t_j = q'_{kj} \times q_{kj} + \lambda_{kj} q_{kj} \quad j = 1, 2, \quad \underbrace{F^k(v)}_{5 \times 3} \begin{bmatrix} \lambda_{k1} \\ \lambda_{k2} \\ 1 \end{bmatrix} = 0$$

- $F^k(v)$ is quadratic in v
- rank condition gives 6th order equations on v

$$\begin{bmatrix} [2] & [2] & [2] \\ [2] & [2] & [2] \\ [2] & [2] & [2] \\ [2] & [2] & [2] \\ [2] & [2] & [2] \end{bmatrix}$$

Using a second asymmetry

- Choosing $U_1 = \begin{bmatrix} 0 & 0 & 0 & 1 \end{bmatrix}^T$ gives 10 linearly independent equations.
- Choosing $U_2 = \begin{bmatrix} 0 & 0 & 0 & 1 \end{bmatrix}^T$ gives 10 linearly independent equations.

Together this gives 14 equations

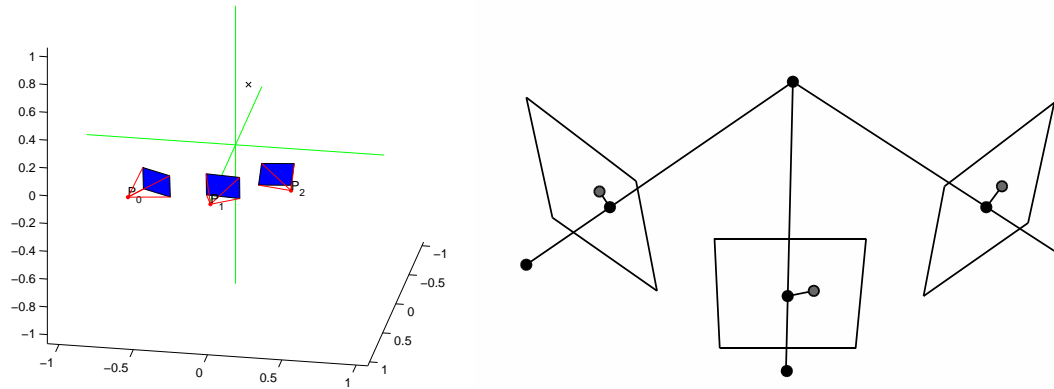
Constructing a GB

1. Add the polynomials $v_j f_i$ there are now 56 polynomials
2. Compute the missing polynomial as an S-polynomial

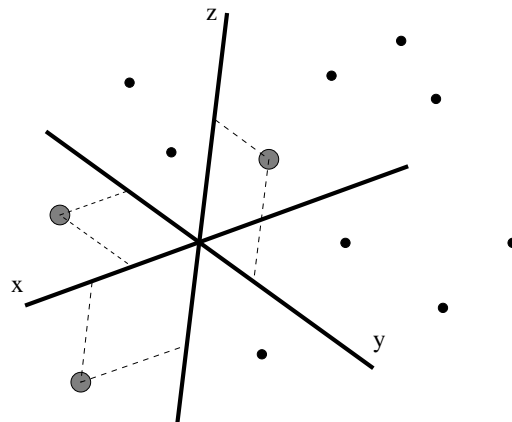
We now have a GrevLex Gröbner-basis.

- As there are 64 inner points to the Gröbner basis the action matrix is 64×64 . This gives a total of 64 solutions.

3-view triangulation, L_2 -norm



- There are 47 stationary points to the target function
- Have to remove unwanted zeros (saturation)



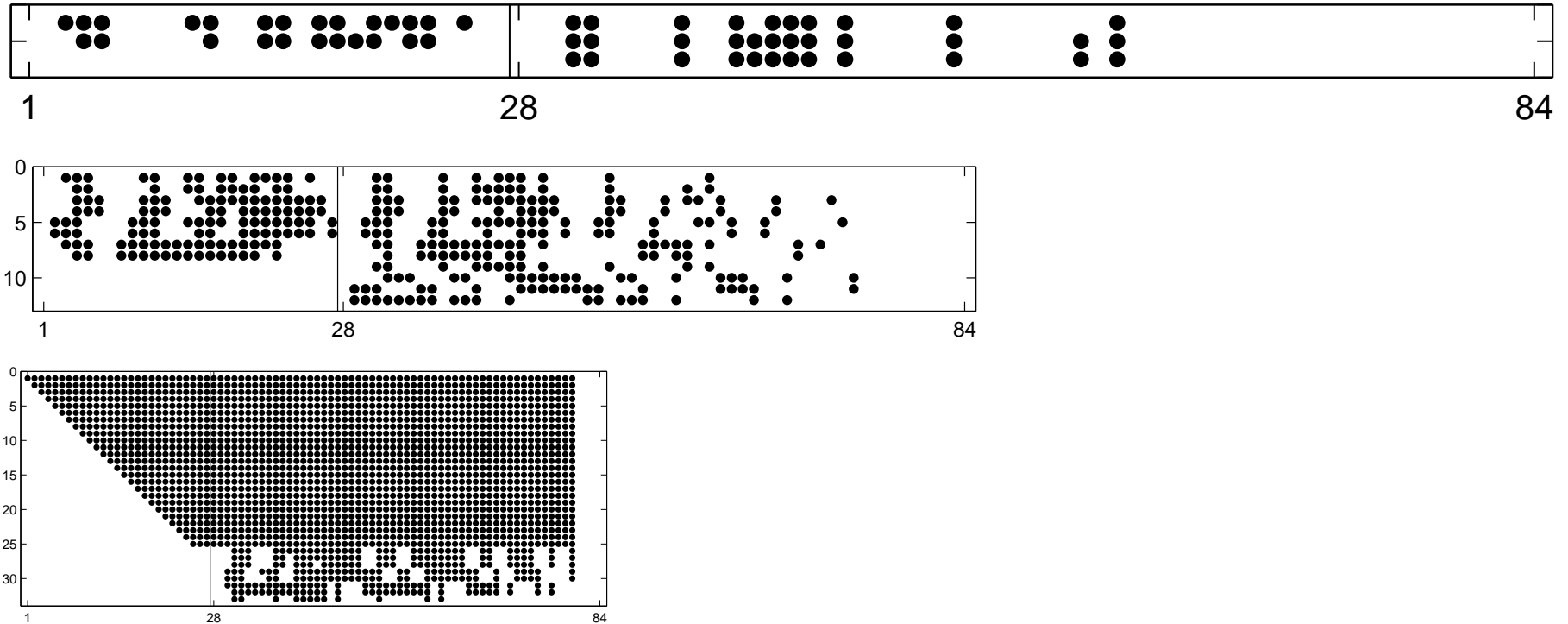
Removing Zeros

We want to remove the set defined by $f(x) = 0$ from the solutions. Two different possibilities:

- Rabinovichs trick. Add the synthetic variable s and the polynomial $1 - sf(x)$.
- Divide out the unwanted solutions.
 - Use special monomial order. (Elimination order)

Saturation

- Saturation, three rounds



- Data flow



Some help

- Construct test to verify results halfway ...
- Do not write by hand what the computer can write for you
- Feel free to browse in my poorly commented code archive
- The slides will be available soon

`http://www.vis.uky.edu/~stewe/`

End

The slides and some code will be available on:

<http://www.vis.uky.edu/~dnister/>

<http://www.vis.uky.edu/~stewe/>